

Interactive Simplifier Tracing and Debugging in Isabelle

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- 1 State of the Art
- 2 Features of the New Simplifier Trace
- 3 Challenges & Open Problems
- 4 Evaluation



- ▶ interactive proof assistant
- ▶ powerful automation
 - ▶ classical and equational reasoning
 - ▶ decision procedures (e.g. linear arithmetic)
 - ▶ integration with external automated theorem provers
 - ▶ ...
- ▶ IDE with continuous proof checking based on jEdit



- ▶ one of the core tactics in Isabelle
- ▶ huge: more than 1800 lines of code
- ▶ applies rewrite rules to terms
- ▶ rules can be conditional: preconditions solved recursively
- ▶ rules can be lazy: “simprocs” can generate rules on the fly
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Simplifier

Example: Conditional rewrite rules



$x, y \in \mathbb{N}$

$$2 \cdot x = x + x \quad (1)$$

$$x < y \implies x - y = 0 \quad (2)$$

$$0 < x + 1 \quad (3)$$

$$0 < x \implies 0 < y \implies 0 < x + y \quad (4)$$

Simplifier



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Simplifier



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Simplification might go wrong:

- ▶ no result at all
- ▶ unexpected result
- ▶ non-termination



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tackled by **tracing**

Simplifier Trace



Lists all rewriting steps, but:

- ▶ potentially huge
- ▶ can't be filtered (e.g. "trace only applications of X and Y")
- ▶ offers no hierarchical structure
- ▶ problematic with non-termination



[1]SIMPLIFIER INVOKED ON THE FOLLOWING TERM:

$10^5 * 10^5 = 10^{10}$

[1]Applying instance of rewrite rule "Power.semiring_numeral_class.power_numeral"
numeral ?k1 ^ numeral ?l1 \equiv numeral (Num.pow ?k1 ?l1)

[1]Rewriting:

$10^5 \equiv$ numeral (Num.pow (num.Bit0 (num.Bit1 (num.Bit0 num.One))) (num.Bit1 (nu

[1]Applying instance of rewrite rule "Num.pow.simps_3":

Num.pow ?x1 (num.Bit1 ?y1) \equiv Num.sqr (Num.pow ?x1 ?y1) * ?x1

[1]Rewriting:

Num.pow (num.Bit0 (num.Bit1 (num.Bit0 num.One))) (num.Bit1 (num.Bit0 num.One)) \equiv
Num.sqr (Num.pow (num.Bit0 (num.Bit1 (num.Bit0 num.One))) (num.Bit0 num.One)) * r

[1]Applying instance of rewrite rule "Num.pow.simps_2":

Num.pow ?x1 (num.Bit0 ?y1) \equiv Num.sqr (Num.pow ?x1 ?y1)

[1]Rewriting:

Num.pow (num.Bit0 (num.Bit1 (num.Bit0 num.One))) (num.Bit0 num.One) \equiv Num.sqr (N

[1]Applying instance of rewrite rule "Num.pow.simps_1":

Num.pow ?y num.One \equiv ?y

[1]Rewriting:

Num.pow (num.Bit0 (num.Bit1 (num.Bit0 num.One))) num.One \equiv num.Bit0 (num.Bit1 (n

Agenda



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Overview



- ▶ interactive
- ▶ breakpoints on terms and theorems
- ▶ configurable verbosity
- ▶ integrated into Isabelle/jEdit



Demonstration

Related Work

SWI-Prolog



- ▶ offers interactive tracing
- ▶ supports breakpoints
- ▶ speciality: marking goals as success

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 - ▶ In Isabelle: difficult because of proof kernel

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Maude



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- ▶ supports breakpoints
- ▶ speciality: during rewriting, issue new goal
 - ▶ In Isabelle: rarely needed because of parallel processing

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Challenges

Selective Memory Clearing



Scenario

1. rewrite step fails
2. user chooses to redo the step
3. simplification starts anew
4. memoization kicks in, step fails again

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Challenges

Context Handling



- ▶ simplification result depends on local assumptions
- ▶ memoization might not make sense across different contexts

$$(P \implies P) \implies (Q \implies P) \implies R$$

Challenges

User Experience



- ▶ user feedback is generally positive
- ▶ used for detecting erratic rules, analyzing simplifier runtime, ...
- ▶ very flexible, but: every additional option generates confusion

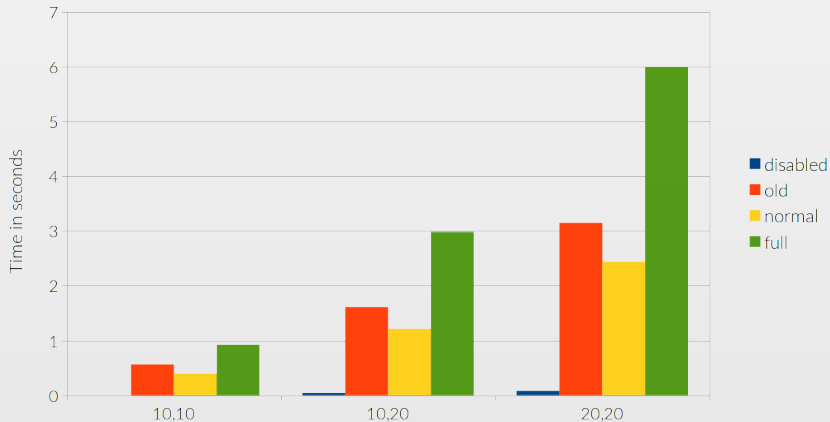
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Performance

Simplifying $10^x \cdot 10^y$



Test machine: Core i7, 3.7 GHz

A Parallelized Simplifier?



- ▶ tracing is completely asynchronous
- ▶ supports multiple questions at the same time
- ▶ but: unused by the simplifier
- ▶ proof of concept: development of a tiny, parallel simplifier

A Parallelized Simplifier?

Lessons Learned



Advantages

- ▶ almost trivial to implement for a toy simplifier
- ▶ GUI part works out of the box



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Disadvantages

- ▶ potentially confusing for users
 - ▶ lots of spurious messages
 - ▶ better filtering required?
 - ▶ holding back messages required?



- ▶ a *generic* tracing facility
 - ▶ using its interface requires little changes to a tactic
 - ▶ parallelization-ready
 - ▶ but not 100% there yet
- ▶ first steps towards instrumenting the simplifier
 - ▶ Should all tactics be written in continuation-passing style?



- ▶ support for more tactics
- ▶ support for other traces (unifier, simp debug, ...)
- ▶ memoization: fuzzy matching
- ▶ term provenance (“Where does that ‘5’ come from?”)
- ▶ tighter integration into Isabelle/jEdit



Q & A